# **Microprocessor Systems**

**Chapter 11** 

**Interrupt Interface of 8088/8086 Microprocessors** 

# Lecture Outline

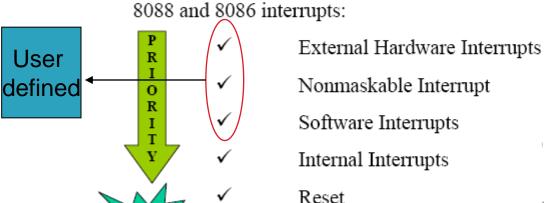
- → 11.1-11.12
- ▶ All sections are included....see page 576

#### INTERRUPT INTERFACE

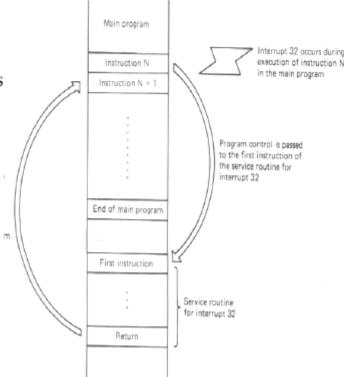
Interrupts provide a mechanism for quickly changing program environment.

The section of the program which the control is passed: Interrupt Service Routine,

ex: For printers it is the printer driver.



256 interrupts are supported by 8088/8086. They are Categorized into 5 Categories



#### 8088/8086 Interrupts

- An interrupt is an external event which informs the CPU that a device needs service
- In the 8088 & 8086 there are a total of 256 interrupts (or interrupt types)
  - INT 00
  - INT 01
  - ...
  - INT FF
- When an interrupt is executed, the microprocessor automatically saves the flags register (FR), the instruction pointer (IP) and the code segment register (CS) on the stack and goes to a fixed memory location.
- In 80x86, the memory location to which an interrupt goes is always four times the value of the interrupt number
- INT 03h goes to 000Ch

#### **Interrupt Service Routine**

- For every interrupt, there must be a program associated with it
- This program is called an Interrupt Service Routine (ISR)
- It is also called an interrupt handler
- When an interrupt occurs, CPU runs the interrupt handler but where is the handler?
  - In the interrupt Vector Table (IVT)

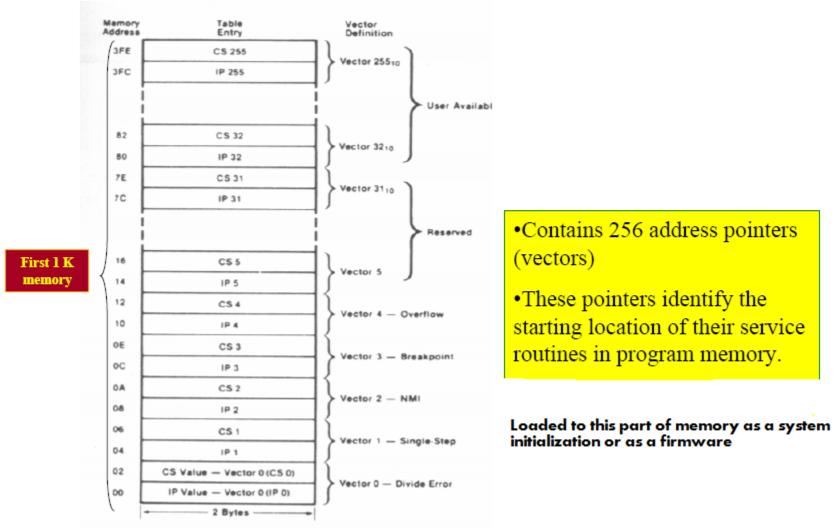
_			
INT Number	Physical Address	Contains	
INT 00	00000h	IP0:CS0	Vector
INT 01	00004h	IP1:CS1	Table
INT 02	00008h	IP2:CS2	
INT FF	003FCh	IP255:CS255	

logical address

#### Interrupt Vector Table

- Interrupt vector table consists of 256 entries each containing 4 bytes.
- Each entry contains the offset and the segment address of the interrupt vector each 2 bytes long.
- Table starts at the memory address 00000H.
- First 32 vectors are spared for various microprocessor operations.
- The rest 224 vectors are user definable.
- The lower the vector number, the higher the priority.

#### Interrupt Vector Table



Simple firmware typically reside in ROM or OTP/PROM, while more complex firmware often employ flash memory to allow for updates. Common reasons for updating firmware include fixing bugs or adding features to the device

#### Example

For example: vector 50: CS and IP?

Physical Address  $200 = (4 \times 50) = 200 = 11001000 = C8H$ 

000C8 contains IP: and 000CA contains CS information

- INT 12h (or vector 12)
- The physical address 30h (4 x 12 = 48 = 30h)
   contains

0030h and 0031h contain IP of the ISR

0032h and 0033h contain CS of the ISR

## Interrupt instructions

- Interrupt enable flag (IF) causes external interrupts to be enabled.
- INT n initiates a vectored call of a subroutine.
- INTO instruction should be used after each arithmethic instruction where there is a possibility of an overflow.
- HLT waits for an interrupt to occur.
- WAIT waits for TEST input to go high.

# Interrupt Instructions

Mnemonic	Meaning	Format	Operation	Flags Affected
CLI	Clear interrupt flag	CLI	0 → (IF)	IF
STI	Set interrupt flag	STI	1 → (IF)	IF
INT n	Type n software interrupt	INT n	(Flags) → ((SP) - 2) 0 → TF, IF (CS) → ((SP) - 4) (2 + 4 · n) → (CS) (IP) → ((SP) - 6) (4 · n) → (IP)	TF, IF
IRET	Interrupt return	IRET	$((SP)) \rightarrow (IP)$ $((SP) + 2) \rightarrow (CS)$ $((SP) + 4) \rightarrow (Flags)$ $(SP) + 6 \rightarrow (SP)$	All
INTO	Interrupt on overflow	INTO	INT 4 steps	TF, IF
HLT	Halt	HLT	Wait for an external interrupt or reset to occur	None
WAIT	Wait	WAIT	Wait for TEST input to go active	None

## The Operation of Real Mode Interrupt

- 1. The contents of the FLAG REGISTERS are pushed onto the stack
- Both the interrupt (IF) and (TF) flags are cleared. This disables the INTR pin and the trap or single-step feature. (Depending on the nature of the interrupt, a programmer can unmask the INTR pin by the STI instruction)
- The contents of the code segment register (CS) is pushed onto the stack.
- 4. The contents of the instruction pointer (IP) is pushed onto the stack.
- The interrupt vector contents are fetched, and then placed into both IP and CS so that the next instruction executes at the interrupt service procedure addressed by the interrupt vector.
- While returning from the interrupt-service routine by the instruction IRET, flags return to their state prior to the interrupt and and operation restarts at the prior IP address.

#### INT 00 (divide error)

```
MOV AL,92
SUB CL, CL
DIV CL ; 92/0 undefined
```

; Also invoked if the quotient is too large to fit into the assigned register

MOV AX,0FFFh MOV BL,2 DIV BL

#### ; WRITE A DIVIDE ERROR ISR

Prompt db 'Division by zero attempted\$'

Divern: PUSH DX Mov ah,09h Mov dx, offset prompt int 21h POP DX

## INT 01 (Single Step)

- **★**In executing a sequence of instructions, there is often a need to examine the contents of the CPU's registers and system memory.
- **★**This is done by executing one instruction at a time and then inspecting the registers and memory
- **★**This is called the tracing or the single stepping
- \*TF must be set (D8 of the flag register)

PUSHF POP AX OR AX,0000000100000000B PUSH AX POPF

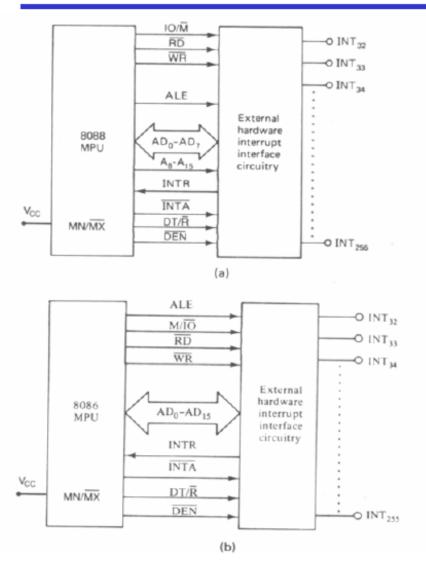
#### Other Interrupts

- INT 02h
  - Intel has set aside INT 02h for the NMI interrupt
  - There is an NMI pin on the CPU
  - If the NMI pin is activated by a H signal, the CPU jumps to 00008H to fetch the CS:IP of the ISR associated with NMI
- INT 03h (breakpoint)
- INT 04H (signed number overflow) or INTO
  - If OF=0 goes to 00010h to get the address of the ISR
  - Otherwise, it is equivalent to NOP
- Example: Use debug dump command to see the IVT
  - D 0000:0000 0013

#### Differences between INT and CALL

- ❖A CALL FAR instruction can jump any location within the 1 MB address range but INT nn goes to a fixed memory location in the Interrupt Vector Table to get the address of the interrupt service routine
- ❖A CALL FAR instruction is used by the programmer in the sequence of instruction in the program but externally activated hardware interrupt can come at any time
- ❖A CALL FAR cannot be masked but INT nn in hardware can be blocked.
- ❖A CALL FAR saves CS:IP but INT nn saves Flags and CS:IP
- ❖At the end of the subroutine RET is used whereas for Interrupt routine IRET should be the last statement

#### **External Hardware Interrupt Interface**



#### Minimum Mode

- ✓ The interrupt circuitry must identify which of the pending interrupts has the highest priority.
- ✓Then passes its type number to the MPU
- ✓ The MPU samples the INTR at the **last clock period** of **each** instruction execution cycle. Its active high level must be maintained.
- ✓ When recognized INTRA generated.

## External hardware-interrupt Interface

- Minimum mode hardware-interrupt interface:
  - 8088 samples INTR input during the last clock period of each instruction execution cycle. INTR is a level triggered input; therefore logic 1 input must be maintained there until it is sampled. Moreover, it must be removed before it is sampled next time. Otherwise, the same Interrupt Service is repeated twice.
  - INTA goes to 0 in the first interrupt bus cycle to acknowledge the interrupt after it was decided to respond to the interrupt.
  - It goes to 0 again the second bus cycle too, to request for the interrupt type number from the external device.
  - The interrupt type number is read by the processor and the corresponding int. CS and IP numbers are again read from the memory.

# HOMEWORK 1

```
FROM THE REVIEW PROBLEMS

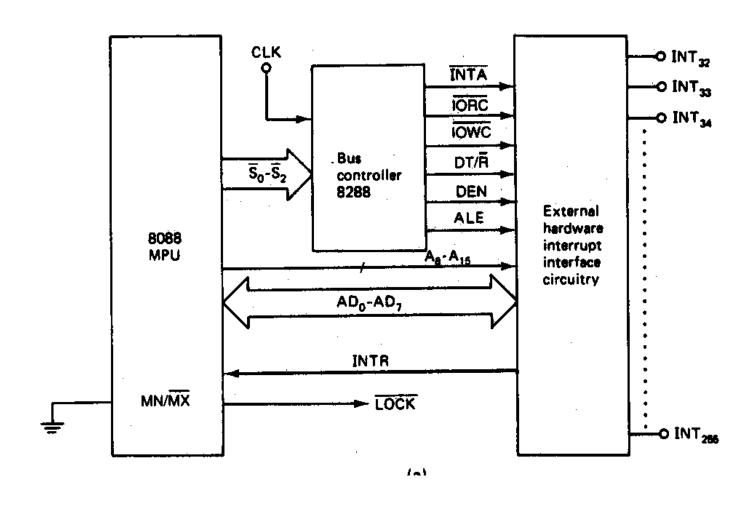
SECTION 11.1: 1, 3, 5

SECTION 11.2: 6, 7, 9, 11

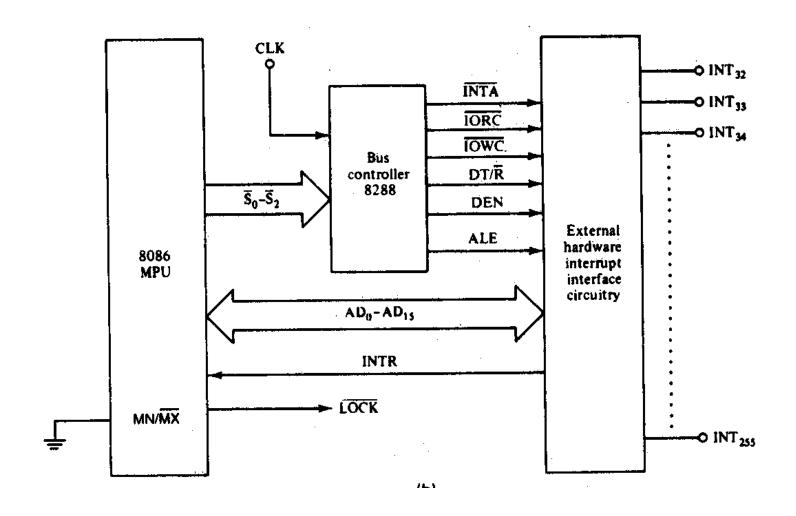
SECTION 11.3: 13

SECTION 11.4: 16
```

## Maximum Mode-External hardware interrupt.

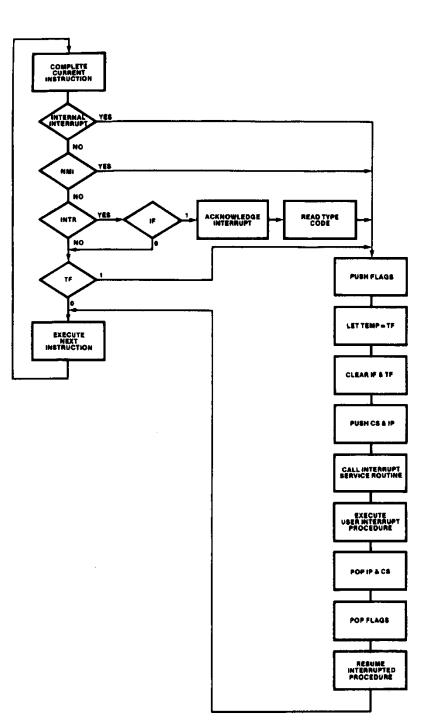


## Maximum Mode-External hardware interrupt (cont'd)



Status inputs			CDU	0000	
$\overline{S}_2$	$\overline{\mathbf{S}}_1$	$\overline{S}_0$	CPU cycle	8288 command	
0	0	0	Interrupt acknowledge	ĪNTĀ	
0	0	1	Read I/O port	ĪŌRC	
0	1	0	Write I/O port	IOWC, AIOWC	
0	. 1	1	Halt	None	
1	0.	0	Instruction fetch	MRDC	
1	0	1	Read memory	MRDC	
1	1	0	Write memory	MWTC, AMWC	
] 1	1	1	Passive	None	

11.6 External
HardwareInterrupt
Sequence



# **External hardware-interrupt Sequence**

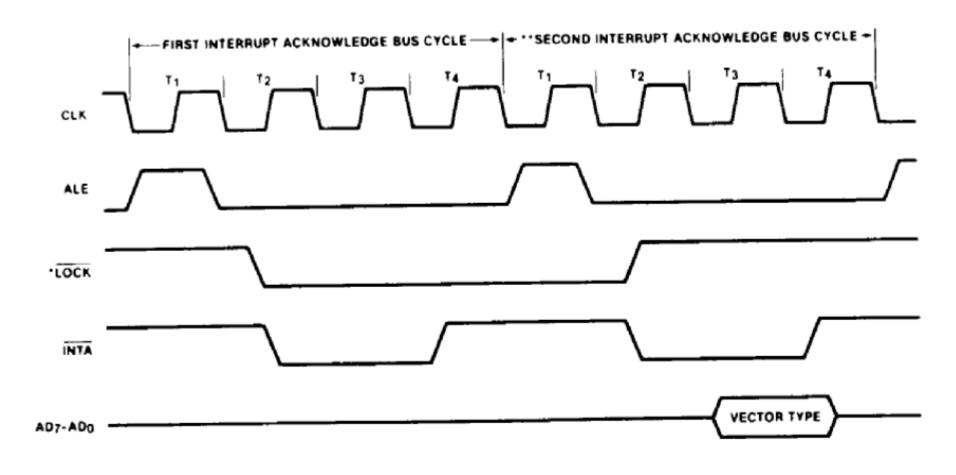
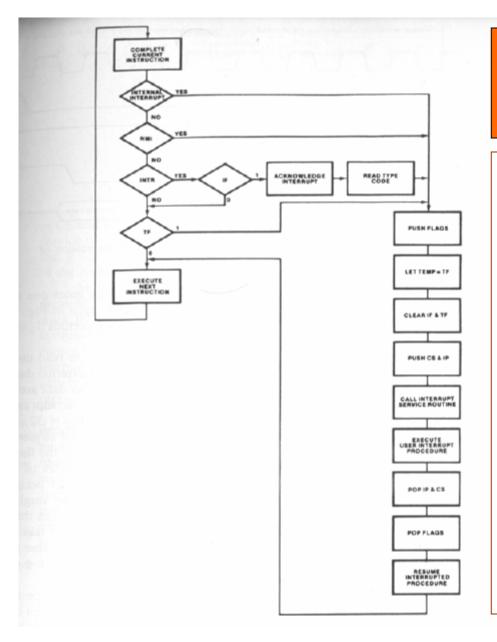


Figure 11-9 Interrupt-acknowledge bus cycle. (Reprinted by permission of Intel Corporation. Copyright/Intel Corp. 1979)



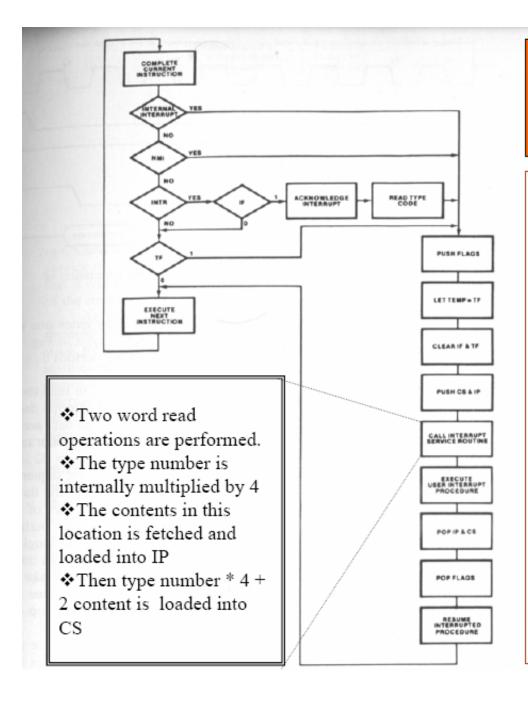
#### Interrupt Sequence

- The interrupt sequence begins when external device requests service by activating one of the interrupt inputs.
- The external device evaluates the priority of this interrupt
- $\triangleright$ INTR  $\rightarrow$  1
- ➤80x86 checks for the INTR at the last T state of the instruction
- Check for IF before granting INTA

# COMPLETE CURRENT INSTRUCTION LET TEMP - TE CLEAR IF & TE PUSH CS & IF POP IP & CS PROCEDURE

#### Interrupt Sequence

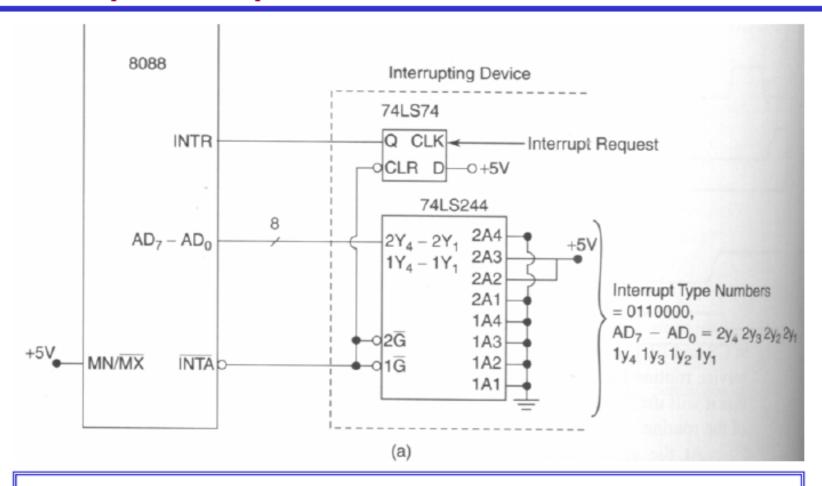
- ➤80x86 initiates the INTA bus cycle. During T1 of the first bus cycle ALE is sent and bus is at Z state and stays high for the bus cycle.
- LOCK is provided in maxmode operation
- During the second interrupt acknowledge bus cycle, external circuitry gates one of the interrupts 20→FF onto data bus lines
- Must be valid during T3 and T4 of second bus cycle



#### Interrupt Sequence

- ➤DT/R and DEN are at logic zero and IO/M is at 1.
- Next save the contents of the flag register
- ➤TF and IF are cleared
- ➤CS and IP are pushed
- ➤Upon return by IRET
- ➤CS and IP are popped
- ➤Flags are popped

## Interrupt Example

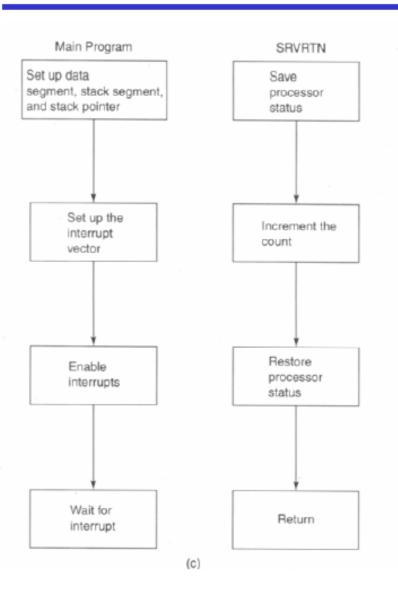


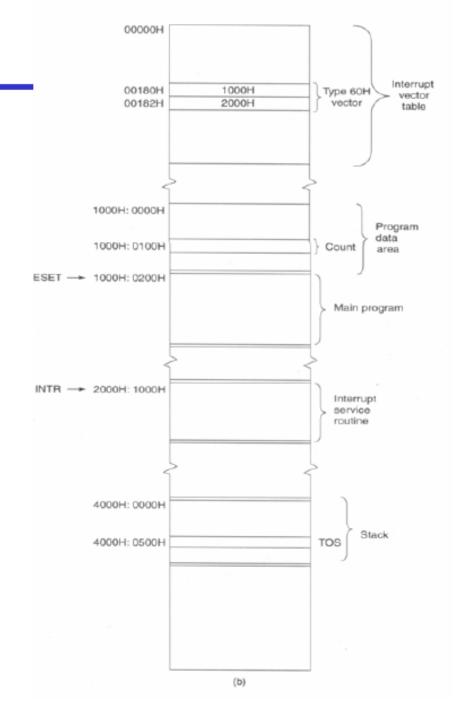
Interrupts the microprocessor each time the interrupt request signal has a transition from 0→ 1. The corresponding interrupt number generated by the hardware in response to INTA is 60H

## Interrupt Example

- An interrupting device interrupts the microprocessor each time the interrupt request input has a transition from 0 to 1.
- 74LS244 creates the interrupt type number 60H as a response to INTA
- Assume:
  - CS=DS=1000H
  - SS=4000H
  - Main program offset is 200H
  - Count (counts the number of interrupts) offset is 100H
  - Interrupt-service routine code segment is 2000H
  - Interrupt-service routine code offset is 1000H
  - Stack has an offset of 500H to the current stack segment
  - Make a map of the memory space organisation
  - Write a main program and a service routine to count the number of positive interrupt transitions.

## Memory organization





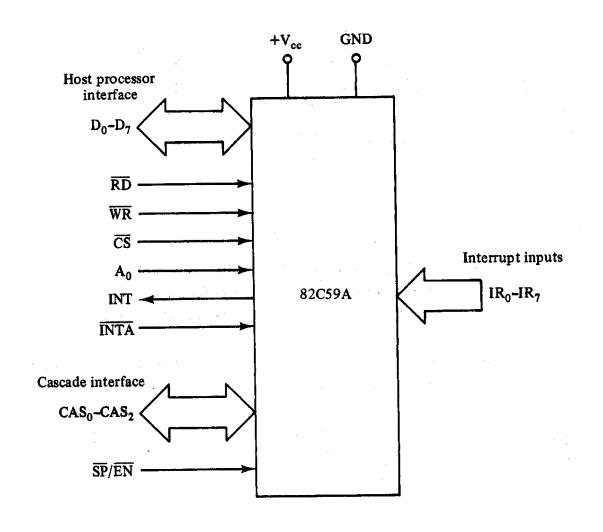
#### **Program**

```
; Main Program, START = 1000H:0200H
START:
           MOV AX, 1000H
                                  ;Setup data segment at 1000H:0000H
           MOV DS, AX
           MOV AX,4000H
                                  ;Setup stack segment at 4000H:0000H
           MOV SS, AX
           MOV SP,0500H
                                  ;TOS is at 4000H;0500H
           MOV AX,0000H
                                  ; Segment for interrupt vector table
           MOV ES, AX
           MOV AX,0000H
                                  :Service routine offset
           MOV [ES:180H], AX
           MOV AX,2000H
                                  ;Service routine segment
           MOV [ES:182H], AX
           STI
                                  ; Enable interrupts
HERE:
           JMP HERE
                                  ; Wait for interrupt
;Interrupt Service Routine, SRVRTN = 2000H:1000H
SRVRTN:
           PUSH AX
                                  ; Save register to be used
           MOV AL, [0100H]
                                  :Get the count
           INC AL
                                  ; Increment the count
           DAA
                                  ; Decimal asdjust the count
           MOV [0100H], AL
                                  ; Save the updated count
           POP AX
                                  ; Restore the register used
           IRET
                                  ; Return from the interrupt
                                   (d)
```

## 8259 Programmable Interrupt Controller

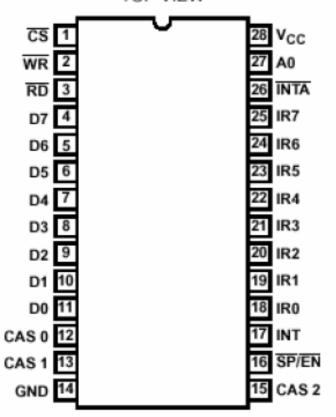
- •The 8259 programmable interrupt controller (PIC) adds eight vectored priority encoded interrupts to the microprocessor.
- •This controller can be expanded to accept up to 64 interrupt requests. This requires a master 8259 and eight 8259 slaves.
- •Vector an Interrupt request anywhere in the memory map.
- •Resolve eight levels of interrupt priorities in a variety of modes, such as <u>fully nested mode</u>, <u>automatic rotation mode</u>, <u>and specific rotation mode</u>.
- •Mask each of the interrupt request individually
- •Read the status of the pending interrupts, in-service interrupts and masked interrupts.

# Block diagram of 82C59



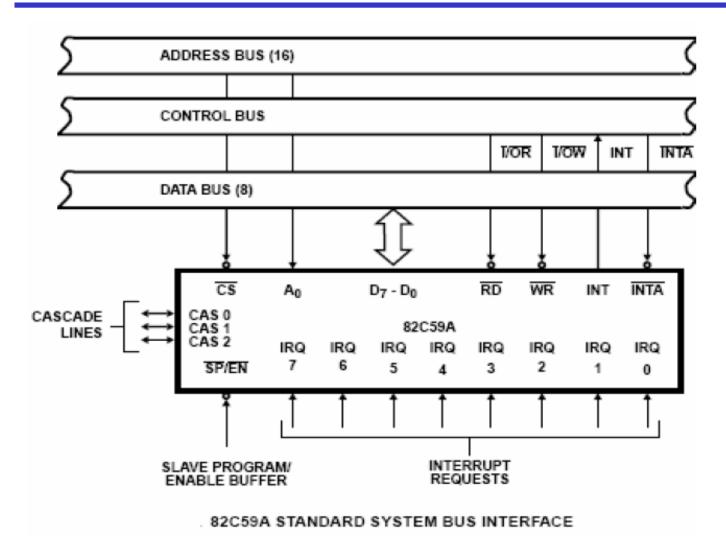
## **Block Diagram**

82C59A (PDIP, CERDIP, SOIC) TOP VIEW



PIN	DESCRIPTION
D7 - D0	Data Bus (Bidirectional)
RD	Read Input
WR	Write Input
A0	Command Select Address
CS	Chip Select
CAS 2 - CAS 0	Cascade Lines
SP/EN	Slave Program Input Enable
INT	Interrupt Output
ĪNTĀ	Interrupt Acknowledge Input
IR0 - IR7	Interrupt Request Inputs

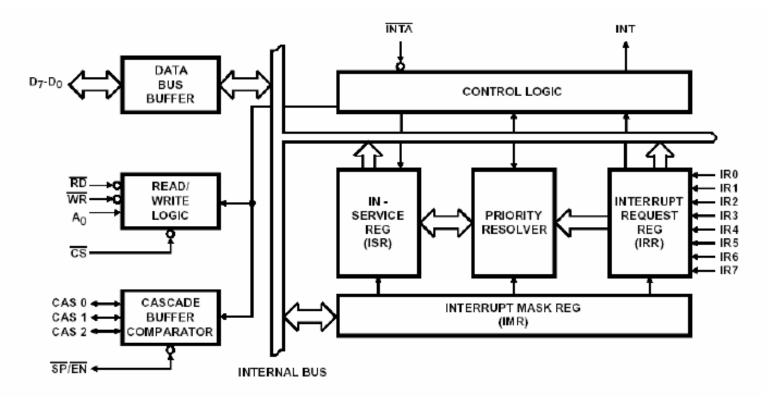
#### 8259 System Bus



34

## 82C59A Programmable Interrupt Controller

- Block diagram of 82C59A includes 8 blocks
  - 8259 is treated by the host processor as a peripheral device.
  - 8259 is configured by the host pocessor to select functions.
- Data bus buffer and read-write logic: are used to configure the internal registers of the chip.
  - A0 address selects <u>different</u> command words within the 8259



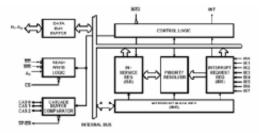
## 82C59A Programmable Interrupt Controller

- Control Logic INT and INTA ared used as the handshaking interface.
  - INT output connects to the INTR pin of the master and is connected to a master IR pin on a slave. INTA is sent as a reply.
  - In a system with master and slaves, <u>only</u> the master INTA signal is connected.
- Interrupt Registers and Priority Resolver: Interrupt inputs IR<sub>0</sub> to IR<sub>7</sub> can be configured as either *level-sensitive* or *edge-triggered* inputs.
   Edge-triggered inputs become active on 0 to 1 transitions.
  - Interrupt request register (IRR): is used to indicate all interrupt levels requesting service.
  - In service register (ISR): is used to store all interrupt levels which are currently being serviced.
  - Interrupt mask register (IMR): is used to enable or mask out the individual interrupt inputs through bits M0 to M7. <u>0= enable, 1= masked out.</u>
  - 4. Priority resolver: This block determines the priorities of the bits set in the IRR. The highest priority is selected and strobed into the corresponding bit of the ISR during the INTA sequence.
    - The priority resolver examines these 3 registers and determines whether INT should be sent to the MPU

# 82C59A Programmable Interrupt Controller

- Cascade-buffer comparator: Sends the address of the selected chip to the slaves in the master mode and decodes the status indicated by the master to find own address to respond.
  - Cascade interface CAS<sub>0</sub>-CAS<sub>2</sub> and SP<sup>-</sup>/EN<sup>-</sup>:
    - Cascade interface CAS<sub>0</sub>-CAS<sub>2</sub> carry the address of the slave to be serviced.
    - SP /EN :=1 selects the chip as the master in cascade mode :=0 selects the chip as the slave in cascade mode :in single mode it becomes the enable output for the data transiver

### Interrupt Sequence

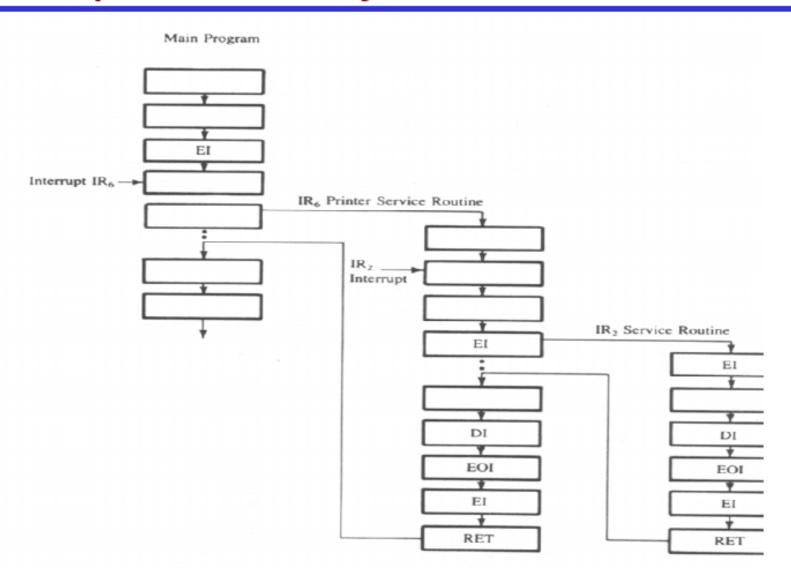


- One or more of the INTERRUPT REQUEST lines (IR0 IR7) are raised high, setting the corresponding IRR bit(s).
- 2) The 82C59A evaluates those requests in the priority resolver with the IMR and ISR, resolves the priority and sends an interrupt (INT) to the CPU, if appropriate.
- The CPU acknowledges the INT and responds with first INTA pulse.
- 4) During this INTA pulse, the appropriate ISR bit is set and the corresponding bit in the IRR is reset (to remove request). The 82C59A does not drive the data bus during the first INTA pulse.
- 5) The 80C86/88/286 CPU will initiate a second INTA pulse. The 82C59A outputs the 8-bit pointer onto the data bus to be read by the CPU.
- 6) This completes the interrupt cycle. In the Automatic End of Interrupt (AEOI) mode, the ISR bit is reset at the end of the second INTA pulse. Otherwise, the ISR bit remains set until an appropriate End of Interrupt (EOI) command is issued at the end of the interrupt subroutine.

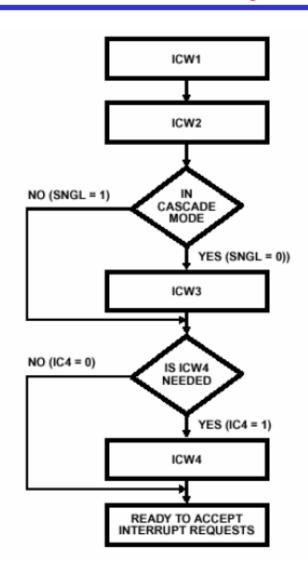
### **Fully Nested Mode**

- It prioritizes the IR inputs such that IR0 has highest priority and IR7 has lowest priority
- This priority structure extends to interrupts <u>currently in service</u> as well as <u>simultaneous interrupt requests</u>
- For example, if an interrupt on IR3 is being serviced (IS3 = 1) and a request occurs on IR2, the controller will issue an interrupt request because IR2 has higher priority.
- But if an IR4 is received (or any interrupt higher than IR2), the controller will not issue the request
- Note however that the IR2 request will not be acknowledged unless the processor has set IF within the IR3 service routine
- In all operating modes, the IS bit corresponding to the active routine must be reset to allow other lower priority interrupts to be acknowledged
- This can be done by outputting <u>manually</u> a special nonspecific EOI instruction to the controller just before IRET
- Alternatively, the controller can be programmed to perform this nonspecific EOI <u>automatically</u> when the second INTA pulse occurs

# **Interrupt Process Fully Nested Mode**



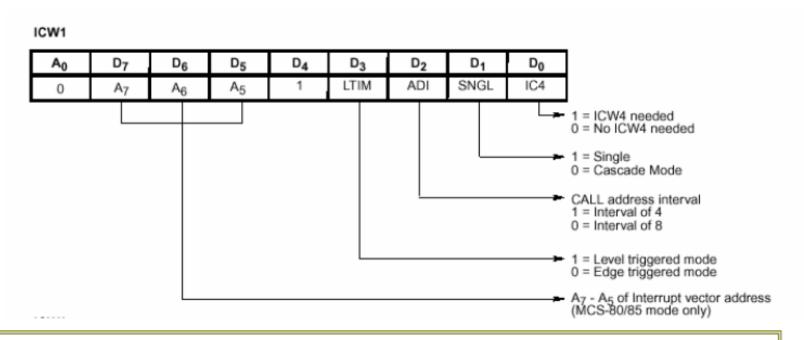
### Initialization Sequence



Two types of command words are provided to program the 8259:

- The initialization command words (ICW)
- The operational command words (OCW)
- Writing ICW1, clears ISR and IMR
- Also Special Masked mode SMM in OCW3, IRR in OCW3 and EOI in OCW2 are cleared to logic 0.
- Fully Nested Mode is entered.
- ICW3 and ICW4 are optional
- It is not possible to modify just one ICW. Whole ICW sequence must be repeated

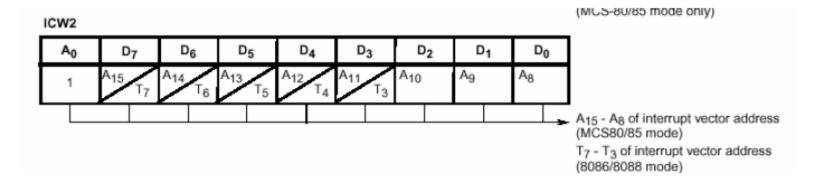
#### ICW1



What value should be written to ICW1 in order to configure the 8259 so that ICW4 needed, the system is going to use <u>multiple</u> 8259s and its inputs are <u>level sensitive</u>?

00011001b = 19h

#### ICW2



What should be programmed into register ICW2 if type number output on the bus is to range from F0h to F7h

$$11110000b = F0h$$

Suppose IR6 is set to generate the value of 6E. Generate the addresses for the other interrupts.

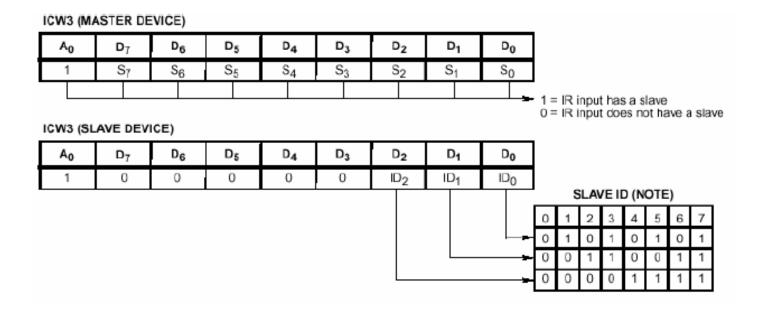
$$IR7 = 6F$$
  $IR3 = 6B$   $IR2 = 6A$   $IR5 = 6D$   $IR1 = 69$   $IR4 = 6C$   $IR0 = 68$ 

# **Content of the Interrupt Vector Byte**

### CONTENT OF INTERRUPT VECTOR BYTE FOR 80C86/88/286 SYSTEM MODE

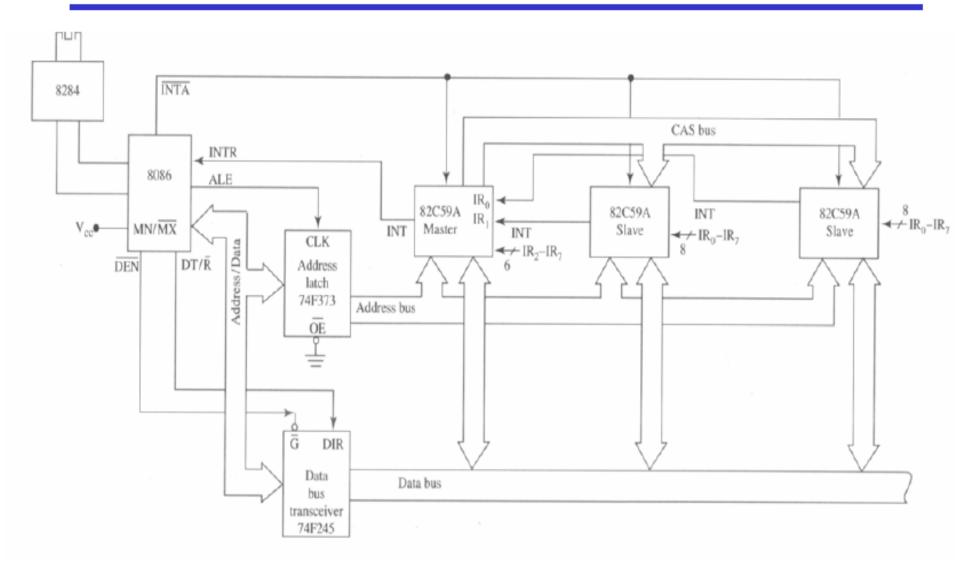
	D7	D6	D5	D4	D3	D2	D1	D0
IR7	T7	T6	T5	T4	T3	1	1	1
IR6	T7	T6	T5	T4	T3	1	1	0
IR5	T7	T6	T5	T4	T3	1	0	1
IR4	T7	T6	T5	T4	T3	1	0	0
IR3	T7	T6	T5	T4	T3	0	1	1
IR2	T7	T6	T5	T4	T3	0	1	0
IR1	T7	T6	T5	T4	T3	0	0	1
IR0	T7	T6	T5	T4	T3	0	0	0

### ICW3

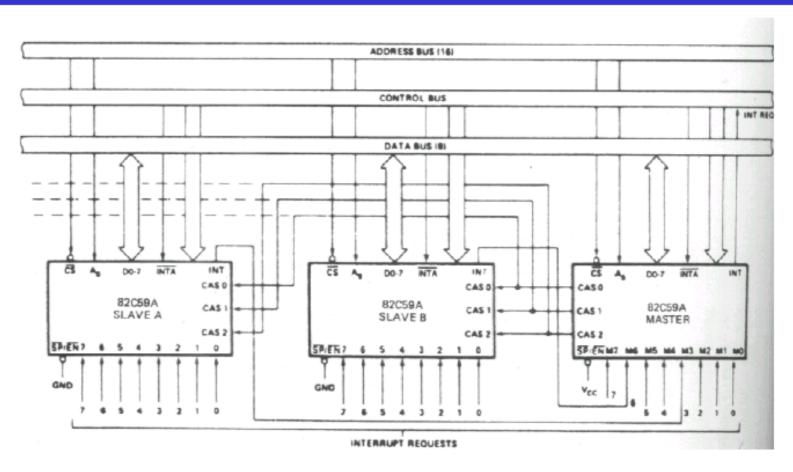


- Q) Suppose we have two slaves connected to a master using IR0 and IR1.
- A) The master is programmed with an ICW3 of 03h, one slave is programmed with an ICW3 of 00h and the other with an ICW3 of 01h.

# Master Slave Configuration

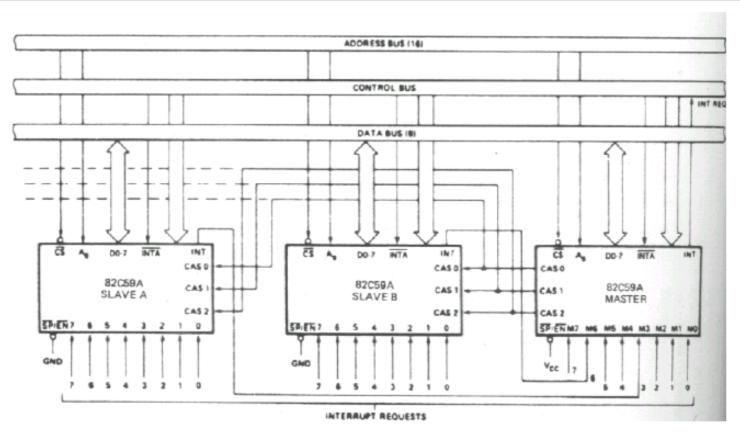


# **Master Slave Configuration**



- ✓ When slave signals the master that an interrupt is active the master determines whether or not its priority is higher than that of any already active interrupt.
- ✓ If the new interrupt is of higher priority the master controller switches INTR to logic 1

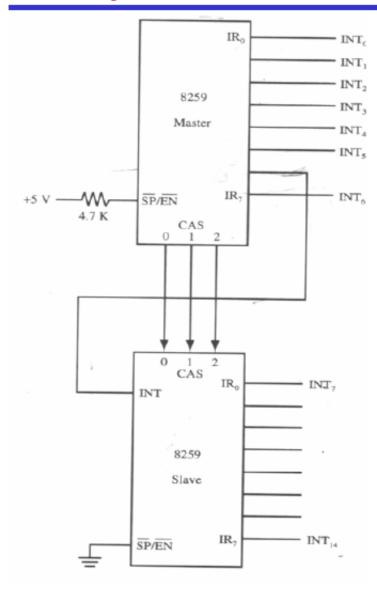
## **Master Slave Configuration**



✓ This signals MPU that external device needs to be serviced. If IF is set. As the first INTA is sent out the master is signaled to output the 3 bit cascade code of the slave device whose interrupt request is being acknowledged on the CAS bus. All slaves read this code and compare internally

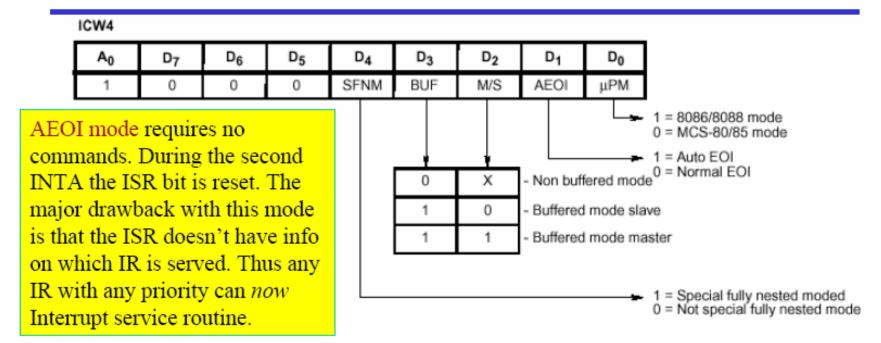
√ The slave corresponding to the code is signaled to output the type number of its highest priority active interrupt on the data bus during the second INTA cycle.

### **Example Master-Slave**



- ✓ Any requests on interrupt lines INT7 through INT14 will cause IR6 to be activated on the MASTER.
- ✓ If so it will output the cascade number of the SLAVE on CAS0 through CAS2.
- ✓ These cascade bits are received by the SLAVE device which examines its ICW3 to see if there is a match...
- ✓ The programmer must have programmed 110 into the SLAVE'S ICW3. If there is a match between the cascade number and ICW3, the SLAVE device will output the appropriate vector number during the second INTA pulse.

#### ICW4



BUF when 1 selects buffer mode. The SP/EN pin becomes an output for the data buffers.

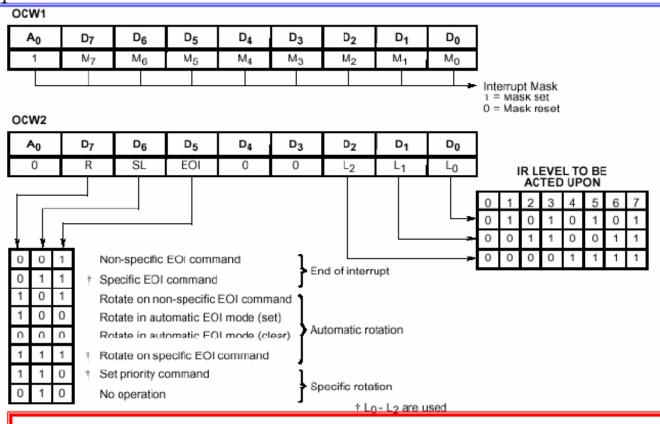
When 0, the SP/EN pin becomes the input for the (MASTER/SLAVE) functionality

M/S is used to set the function of the 8259 when operated in buffered mode. If M/S is set the 8259 will function as the MASTER.

If cleared will function as SLAVE.

#### OCW1 - OCW2

OCW1 is used to access the contents of the IMR. A READ operation can be performed to the IMR to determine the present setting of the mask. Write operations can be performed to mask or unmask certain bits.



Controller will not confuse OCW2 with ICW1 since D4 = 1

### Example

```
ISR PROC FAR
...
MOV AL, 00100000b
OUT 20h, AL
IRET
ISR ENDP
```

What should be OCW1 if interrupt inputs IR0 through IR3 are to be masked and IR4 through IR7 are to be unmasked?

```
D3D2D1D0 = 1111
D7..D4 = 0

→ 00001111 = 0F
```

What should be OCW2; if priority scheme rotate on non specific EOI issued 101 00000 (since it doesn't have to be specific on certain bit

### OCW3

Permits reading of the contents of the ISR or IRR registers through software

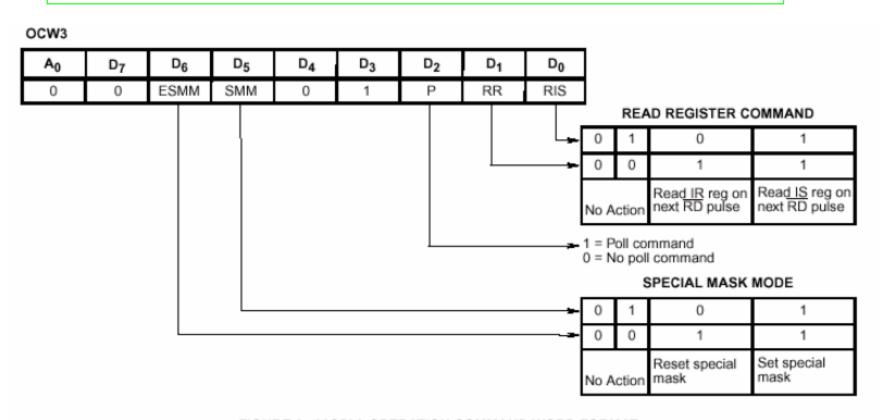


FIGURE 8. 82C59A OPERATION COMMAND WORD FORMAT

### Example

Normally when an IR is acknowledged and EOI is not issued, lower priority interrupts will be inhibited.

So the SPECIAL MASK MODE, when a mask bit is set in OCW1, it inhibits further interrupts at that level and end enables from all other levels, that are not masked.

MOV AL, 00010000b ; mask IRQ4

OUT 21h, AL ; OCW1 (IMR)

MOV AL, 01101000b ; special mask mode

OUT 20h, AL ; OCW3

; by masking itself and selecting the special mask mode interrupts on IRQ5 thru IRQ7 will now be accepted by the controller as well as IRQ0 thru IRQ3

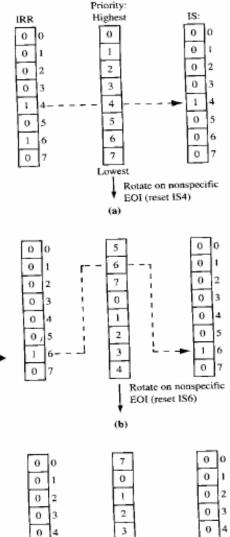
### **End of Interrupt**

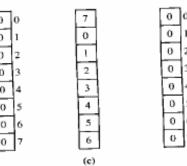
- ➤ The In Service (IS) bit can be reset automatically following the trailing edge of the last in sequence INTA pulse (when AEOI bit in ICW4 is 1) or by a command word that must be issued to the 8259 before returning from a service routine (EOI command).
- An EOI command must be issued **twice** in the Cascade mode, once for the master and once for the corresponding slave.
- ➤ There are two forms of (non-automatic) EOI command:
  - ✓ Specific: When there is a mode which may disturb the fully nested structure, the 8259 may not determine the last level acknowledged. In this case a specific EOI must be issued, which includes the IS level to be reset. (OCW2)
  - ✓ Non Specific: When a Non Specific EOI issued the 8259 will automatically reset the highest IS bit of those that are set, since in the fully nested mode the highest level was necessarily the last level acknowledged and serviced. (preserve the nested structure)
    - ❖ A non Specific EOI can be also issued at OCW2.

### **Automatic Rotation**

interrupt requests arrive on IR4 and IR6

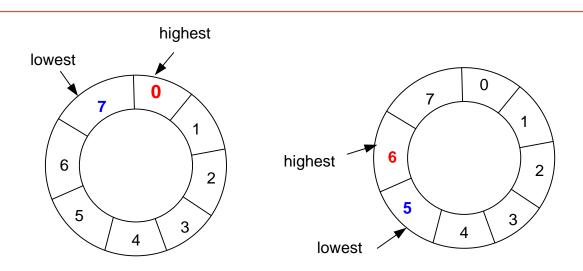
EOI command always resets the highest ISR bit (bit of highest priority) Use automatic rotating mode to clear the IS bit as soon as it is acknowledged





### **Specific Rotation**

- The programmer can change priorities by programming the bottom priority and thus fixing all other priorities (for ex: if IR5 is programmed as the bottom priority device, then IR6 will have the highest one)
- The set priority command is issued in OCW2 where R=1, SL=1, L0-L2 is the binary priority level code of the bottom priority device)



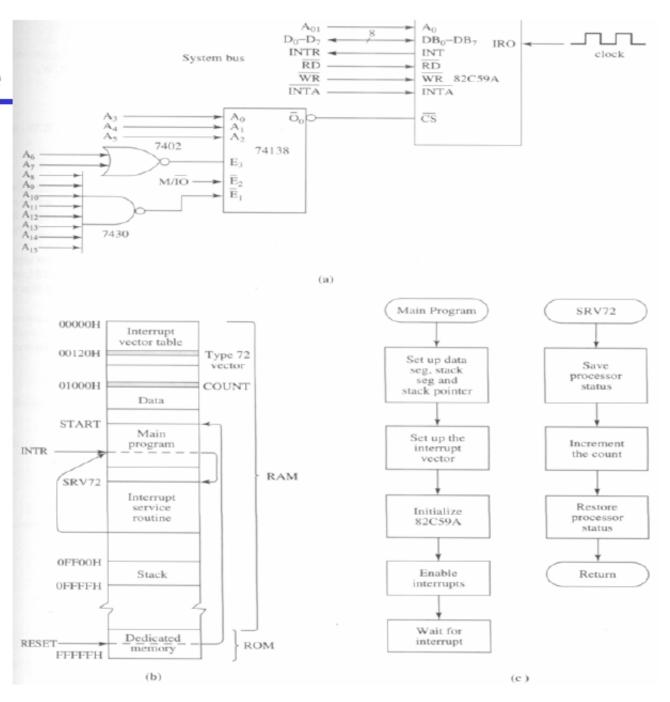
### **Special Fully Nested Mode**

- Used in the case of a large system where cascading is used, and the priority has to be conserved within each slave.
- This mode is similar to the normal nested mode with the following:
  - When an interrupt request from a certain slave is in service this slave is not locked out from the master's priority logic and further interrupt requests from higher priority IR's within the slave will be recognized by the master and will initiate interrupts to the processor.
  - When exiting the ISR the software has to check whether the interrupt is the only interrupt that is serviced from the SLAVE. This is done by sending an EOI command and check the In service register in the SLAVE. If it is the only one, a non specific EOI bas to be sent to the MASTER, if it is not empty no action performed.

### Example

Analyze the circuit and write an appropriate main program and a service routine that counts as a decimal number the positive edges of the clock signal applied to IR0

Use type number 72



# Example (cont'd)

See chip select CS and A0 of the 8259

- A0 not used
- Two I/O addresses are FF00h and FF02h
- FF00h: ICW1,
- FF02h: ICW2, ICW3, ICW4, OCW1
- ICW1 = 00010011b = 13h
- type number 72 will be used
  - ICW2 = 01001000b = 48h
- ICW3 not needed
- nonbuffered and auto EOI
  - ICW4 = 03h
- mask all other interrupts but IR0
  - OCW1 = 111111110b = FEh

# Main program and ISR

```
CLI
```

START: MOV AX, 0

MOV ES, AX

MOV AX, 100h

MOV DS, AX

MOV AX, 0FF0h; stack

MOV SS, AX

MOV SP, 100h

; interrupt install

MOV AX, OFFSET SRV72

MOV [ES:120h], AX

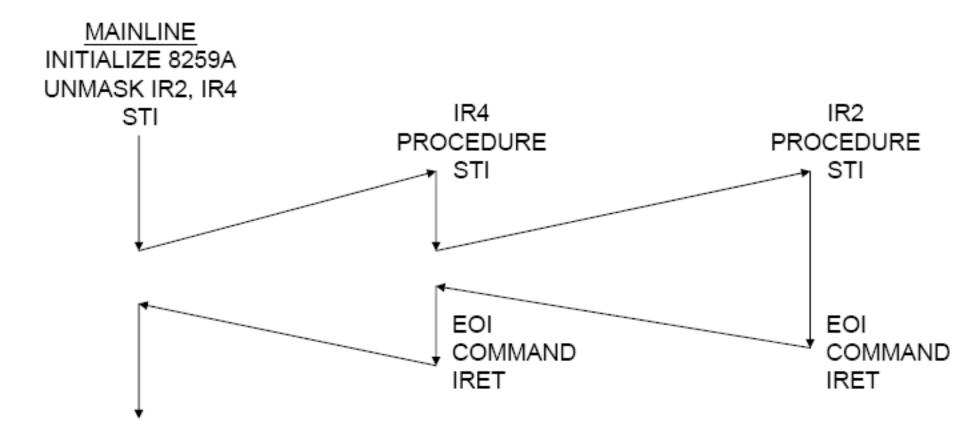
MOV AX, SEG SRV72

MOV [ES:122h]. AX

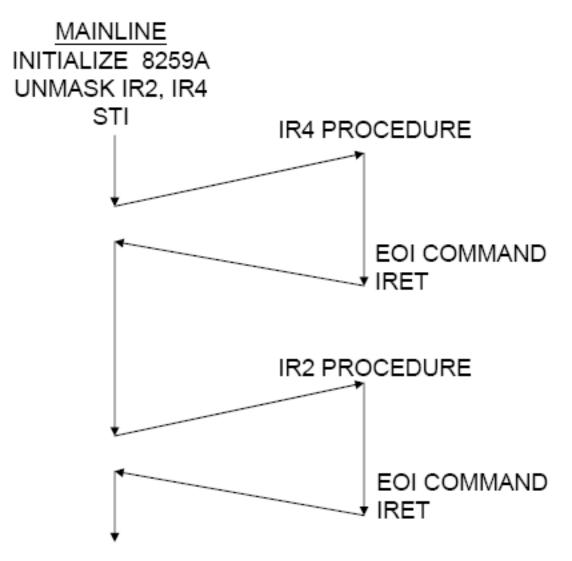
## Example contd

```
; initialization
                              : service routine
      MOV DX, 0FF00h
                               SRV72: PUSH AX
      MOV AL, 13h
                                        MOV AL, [COUNT]
      OUT DX, AL
                                        INC AL
      MOV DX, 0FF02h
                                        DAA
      MOV AL, 48h
                                        MOV [COUNT], AL
      OUT DX, AL
                                        POP AX
      MOV AL, 03h
                                        IRET
      OUT DX, AL
       MOV AL, 0FEh
      OUT DX, AL
       STI
; wait for interrupt
HERE: JMP HERE
```

### IR4 followed by IR2



- (a) Response with INTR enabled inside IR4 procedure (IF=1)
  - Q: what is the default value of I flag in an Interrupt Service Routine?



(b) Response with INTR <u>not</u> enabled inside IR4 procedure (IF=0)

# Homework 2

- Chapter 11
- 27, 30, 32, 34, 36, 37, 40, 51